Impact of Decentralized Learning on Player Utilities in Stackelberg Games Meena Jagadeesan (UC Berkeley)

Joint work with Kate Donahue (Cornell), Nicole Immorlica, Brendan Lucier, Alex Slivkins (MSR)









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High-level overview of this work

We study Stackelberg games with decentralized learning.

Captures systems of two sequential, misaligned agents that learn over time

How do the learning dynamics behave?

Outline for the talk

- 1. Motivation and conceptual overview
- 2. Model of Stackelberg games with decentralized learning
- 3. Our analysis of learning dynamics











Al agent learns from repeatedly interacting with a user.



Al agent learns from repeatedly interacting with a user. User learns from repeatedly interacting with the Al agent.



Al agent learns from repeatedly interacting with a user. User learns from repeatedly interacting with the Al agent.

User-Al interactions are a form of multi-agent learning.







Sequential: User leads (w/ prompt)



Sequential: User leads (w/ prompt); chatbot follows (w/ output).



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Misaligned: User utility (individual preferences) vs. Chatbot utility (e.g., societal preferences, implicit objective learned during training)



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Decentralized: User + chatbot learn separately over a chat session.

We study Stackelberg games with decentralized learning, which are:

- **Sequential**: One agent leads, and the other agent follows.
- **Misaligned:** Agents have different utility functions.
- **Decentralized:** Agents learn separately from repeated interactions.

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How do the learning dynamics behave?

- Implications for each agent's cumulative utility over time?
- Implications for **learning algorithm design**?

Overview of our contributions

Misalignment in agent utilities distorts the learning dynamics.

- When agents can be arbitrarily misaligned, we show the (full-information) Stackelberg equilibrium utilities are **unachievable**.
- We develop **relaxed benchmarks** for each agent's utility, and construct algorithms that perform well for both agents w.r.t. these benchmarks.
- When agents are partially aligned, we show the Stackelberg equilibrium utilities can be achieved.

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- □ **Misaligned** (agent utilities are not equal)
- □ **Decentralized** (agents learn separately)

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Stackelberg games: sequential, misaligned, static environments

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Our model: Stackelberg games with decentralized learning

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Stackelberg games: sequential, misaligned, static environments

Our model: Stackelberg games with decentralized learning

- Agents interact over T rounds.
- At every round, the leader goes first, and the follower goes second.
- Each agent observes their own stochastic rewards.

Recap of static Stackelberg games

	Action	Utility
Leader	$a \in A$	$u_1(a,b)$
Follower	$b \in B$	$u_2(a,b)$

Follower best-responds to leader: $b^*(a) = argmax_{b\in B}u_2(a, b)$

Leader **anticipates follower's actions** and best-responds: $a^* = argmax_{a \in A}u_1(a, b^*(a))$

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Requires full knowledge of u_2

Leader **anticipates follower's actions** and best-responds:

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Requires full knowledge of u_1 and b^*

Each agent learns how to select actions over T rounds.

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We cast agent learning within the **stochastic multi-armed bandit framework**:

Leader

A = arms, u_1 = mean reward function *ALG*₁ = bandit algorithm

Follower

B = arms, u_2 = mean reward function *ALG*₂ = bandit algorithm

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LeaderFollower $A = arms, u_1 = mean reward function$ $B = arms, u_2 = mean reward function$ $ALG_1 = bandit algorithm$ $ALG_2 = bandit algorithm$

At each time step *t*:

Chooses action *a*_t using *ALG*₁

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At each time step <i>t</i> :		
Chooses action a_t using ALG_1	Observes a_t & chooses b_t using ALG_2	
Observe stochastic reward $u_1(a_t, b_t) + \eta_{1,t}$	Observe stochastic reward $u_2(a_t, b_t) + \eta_{2,t}$	

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Cumulative utility: $\sum_{t=1}^{T} u_1(a_t, b_t)$	Cumulative utility: $\sum_{t=1}^{T} u_2(a_t, b_t)$	

Related work

Learning in Stackelberg games with unknown utilities for both agents:

e.g., Camara, Hartline, Johnsen (2020), Bai, Jin, Wang, Xiong (2021), Gan, Han, Wu, Xu (2023), Haghtalab, Podimata, Yang (2023), Collina, Roth, Shao (2023), etc.

Broader literature on learning in Stackelberg games:

e.g., Letchford, Conitzer, Munagala (2009), Balcan, Blum, Haghtalab, Procaccia (2015), Braverman, Mao, Schneider, Weinberg (2018), Fiez, Chasnov, Ratliff (2019), Deng, Schneider, Sivan (2019), Zrnic, Mazumdar, Sastry, Jordan (2021), Kao, Wei, Subramanian (2022), Goktas, Zhao, Greenwald (2022), Haghtalab, Lykouris, Nietert, Wei (2022), Zhao, Zhu, Jiao, Jordan (2023), Brown, Schneider, Vodrahalli (2023), Guruganesh, Kolumbus, Schneider, Talgam-Cohen, Vasileios-Vlatakis-Gkaragkounis (2024), etc.

Interacting learners:

e.g. Chayes, Immorlica, Jain, Etesami, Mahdian (2007), Chan, Hadfield-Menell, Srinivasa, Dragan (2010), Daskalakis, Deckelbaum, Kim (2011), Borgs, Agarwal, Luo, Neyshabur, Schapire (2017), Aridor, Mansour, Slivkins, Wu (2020), Zhuang, Hadfield-Menell (2020), J., Jordan, Haghtalab (2023), etc.

Our model: stochastic rewards, utility of both agents, arbitrary misalignment, decentralized learning

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Measuring each agent's regret

We study each agent i's expected pseudo-regret: $E[\sum_{t=1}^{T} u_i(a_t, b_t)] - \alpha_i \cdot T$ Agent i's cumulative utility Agent i's benchmark

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What benchmarks are appropriate for this environment?

Naïve benchmark: $\alpha_i^{orig} \coloneqq u_i(a^*, b^*(a^*))$ (the Stackelberg equilibrium utility)

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	b_1	. b ₂
<i>a</i> ₁	$(0.6, \delta)$	(0.2, 0)
<i>a</i> ₂	(0.5, 0.6)	(0.4, 0.4)

	b_1	. b ₂
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<i>a</i> ₁	$(0.6, \delta)$	(0.2, 0)	<i>a</i> ₁	$(0.6, \delta)$	(0.2, 2 δ)
a_2	(0.5, 0.6)	(0.4, 0.4)	a ₂	(0.5, 0.6)	(0.4, 0.4)

Pair = $(u_1(a, b), u_2(a, b))$, Gold = follower's best-response Purple = leader's best-response

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Due to misalignment, small errors by one agent can distort the other agent's utility.

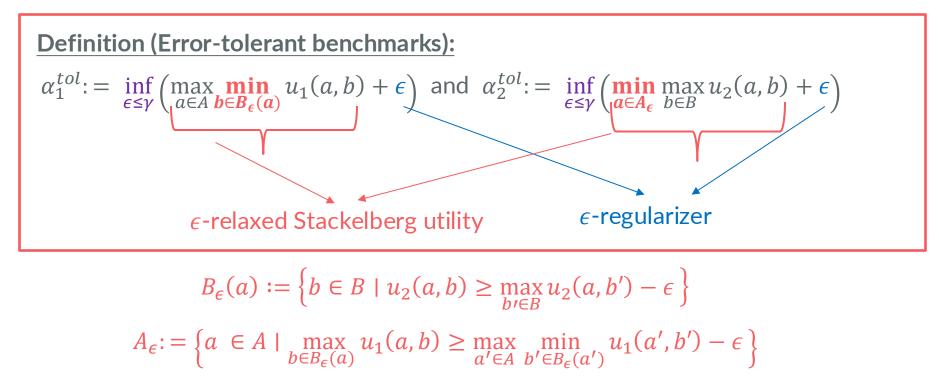
Account for agent errors via ϵ -approximate best-response sets:

Definition (Error-tolerant benchmarks): $\alpha_1^{tol} := \inf_{\epsilon \leq \nu} \left(\max_{a \in A} \min_{b \in B_\epsilon(a)} u_1(a, b) + \epsilon \right) \text{ and } \alpha_2^{tol} := \inf_{\epsilon \leq \nu} \left(\min_{a \in A_\epsilon} \max_{b \in B} u_2(a, b) + \epsilon \right)$ $B_{\epsilon}(a) := \left\{ b \in B \mid u_2(a, b) \ge \max_{b \neq B} u_2(a, b') - \epsilon \right\}$ $A_{\epsilon} := \left\{ a \in A \mid \max_{b \in B_{\epsilon}(a)} u_1(a, b) \ge \max_{a' \in A} \min_{b' \in B_{\epsilon}(a')} u_1(a', b') - \epsilon \right\}$

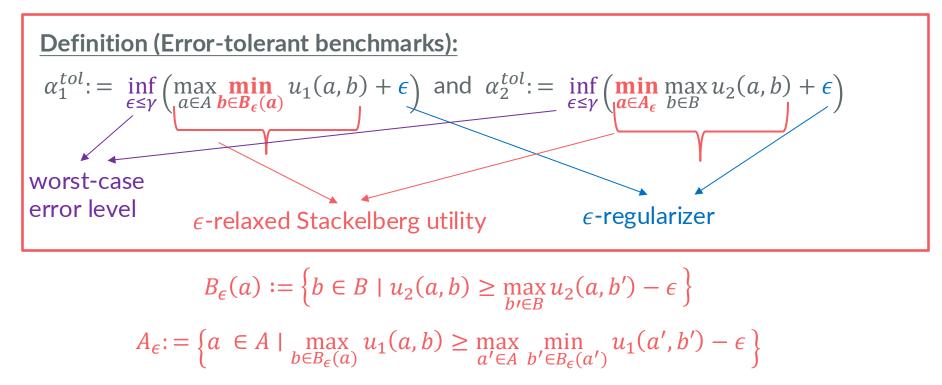
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Definition (Error-tolerant benchmarks): $\alpha_1^{tol} := \inf_{\epsilon \leq \gamma} \left(\max_{a \in A} \min_{b \in B_{\epsilon}(a)} u_1(a, b) + \epsilon \right) \text{ and } \alpha_2^{tol} := \inf_{\epsilon \leq \gamma} \left(\min_{a \in A_{\epsilon}} \max_{b \in B} u_2(a, b) + \epsilon \right)$ ϵ -relaxed Stackelberg utility $B_{\epsilon}(a) := \left\{ b \in B \mid u_2(a, b) \ge \max_{b' \in B} u_2(a, b') - \epsilon \right\}$ $A_{\epsilon} := \left\{ a \in A \mid \max_{b \in B_{\epsilon}(a)} u_1(a, b) \ge \max_{a' \in A} \min_{b' \in B_{\epsilon}(a')} u_1(a', b') - \epsilon \right\}$

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Example of error-tolerant benchmarks

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We take $\delta < \gamma = 0.05$ in this example.

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		Green is $B_{\delta}(a)$	<mark>a)</mark> Pi	nk is A_{δ}		
$(\alpha_1^{tol}, \alpha_2^{tol}) = (0.5 + \delta, \delta)$					$(\alpha_1^{tol}, \alpha_2^{tol}) =$	(0 . 5 , 3 δ)

We take $\delta < \gamma = 0.05$ in this example.

Our algorithmic results

Goal: design algorithms that achieve sublinear regret for both agents

- Standard algorithms can incur linear regret w.r.t the error-tolerant benchmarks.
- We construct algorithms achieving $\tilde{O}(T^{\frac{2}{3}})$ regret w.r.t error-tolerant benchmarks.
- Any algorithms incur $\Omega(T^{\frac{2}{3}})$ regret for some agent w.r.t error-tolerant benchmarks.

• When agent utilities are partially aligned, we construct algorithms which achieve $\tilde{O}(T^{\frac{1}{2}})$ regret w.r.t the **original Stackelberg benchmarks**.



Proposition (Informal): Suppose both agents run ExploreThenCommit. Then **both agents** can **incur linear regret** w.r.t. the error-tolerant benchmarks.

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Leader's estimated rewards during exploration:

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$$0.5(u_1(a_1, b_1) + u_1(a_1, b_2)) = 0.4 \text{ on } a_1$$

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Key issue: the follower's exploration phase distorts the leader's learning

Warm-up: Modified ETC yields sublinear regret

Key algorithmic idea: leader waits for the follower to finish exploring before learning

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Proposition (Informal): Suppose that:

- The follower runs ExploreThenCommit.
- The leader runs a modified ExploreThenCommit where they discard observations from the follower's exploration phase when computing the empirical means.
 => Both agents achieve Õ (T²/₃ (|A||B| log T)¹/₃) regret w.r.t. error-tolerant benchmarks.

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Regret is sublinear for both players!

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Theorem (Informal): Suppose that:

- The leader runs **ExploreThenUCB** where **they discard observations from an initial phase and then run a variant of UCB**.
- The follower runs any algorithm with sufficiently low instantaneous regret. => Both agents achieve $\tilde{O}\left(T^{\frac{2}{3}}\left(|A||B|\log T\right)^{\frac{1}{3}}\right)$ regret w.r.t. error-tolerant benchmarks.

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Follower must gracefully learn (satisfied by AAE, ExploreThenCommit, etc)

Regret scaling with $T^{2/3}$ **rate is** <u>unavoidable</u>

Theorem (Informal): For any ALG₁ and ALG₂, some agent incurs $\Omega\left(T^{\frac{2}{3}}|B|^{\frac{1}{3}}\right)$ regret.



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<i>a</i> ₁	$(0.5 + \delta, \delta)$	(0, 0)
<i>a</i> ₂	$(0.5, 3\delta)$	(0.5 <i>,</i> 3δ)

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To distinguish instances, need to explore a very suboptimal arm (a_1, b_2) for the leader

What if agents are partially aligned?

Suppose that the two agents agree over which outcomes are different.

$$L = \sup_{a,a',b,b'} \left(\frac{u_1(a,b) - u_1(a',b')}{u_2(a,b) - u_2(a',b')}, \frac{u_2(a,b) - u_2(a',b')}{u_1(a,b) - u_1(a',b')} \right)$$

We show that L is bounded => the original Stackelberg benchmarks are achievable.

Theorem (Informal): There exist algorithms such that both players achieve $\widetilde{O}\left(L^2\sqrt{T|A||B|}\right)$ regret w.r.t. the original Stackelberg benchmarks.

Takeaway: Partial alignment makes learning easier and faster.

Conclusion

We study Stackelberg games with decentralized learning.

Main finding: misalignment in agent utilities distorts learning dynamics

- We showed that the Stackelberg equilibrium utilities are unachievable.
- We designed error-tolerant benchmarks to better capture learning dynamics.
- We constructed algorithms which achieve optimal $\tilde{\Theta}(T^{\frac{2}{3}})$ regret.
- We showed that partial alignment makes learning easier and faster.

Future directions: allow for greater flexibility in leader algorithm, study application-specific learning algorithms, characterize equilibria in the meta-game between agents, etc.